

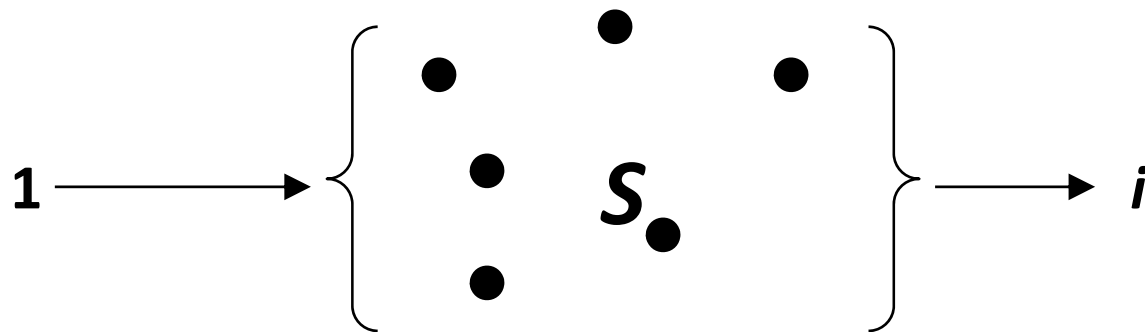
Algorithm Graph Theory: How hard is your combinatorial optimization problem?

Short Course – Lecture 12
June 14, 2017

Slides available at:
<https://www.math2.rwth-aachen.de/de/mitarbeiter/koster/agtclemsun>

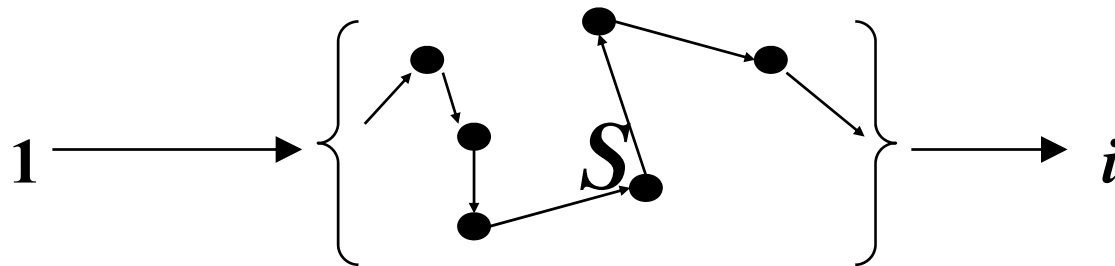
TSP: Dynamic Programming

For any $S \subseteq \{2, 3, \dots, n\}$ and $i \notin S$, let **(S,i)-path** be a path which starts at city 1, visits each city in S exactly once, and no other city, and ends up in city i .



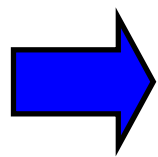
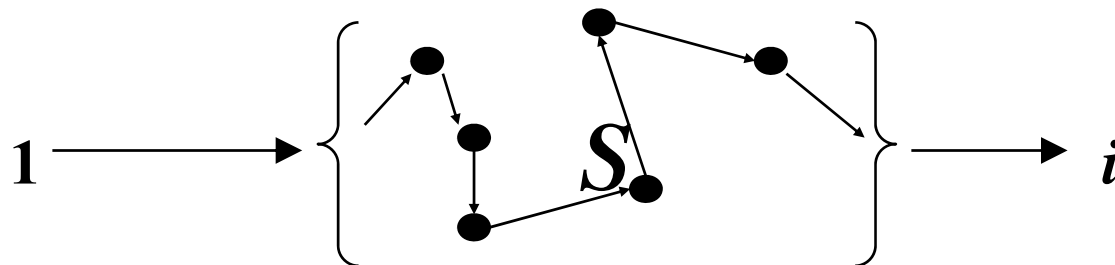
TSP: Dynamic Programming

Let $\mathbf{cost}[S,i]$ stand for the length of the shortest (S,i) path

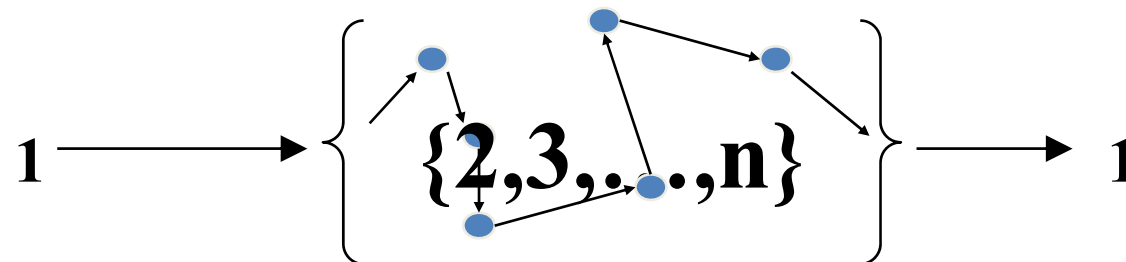


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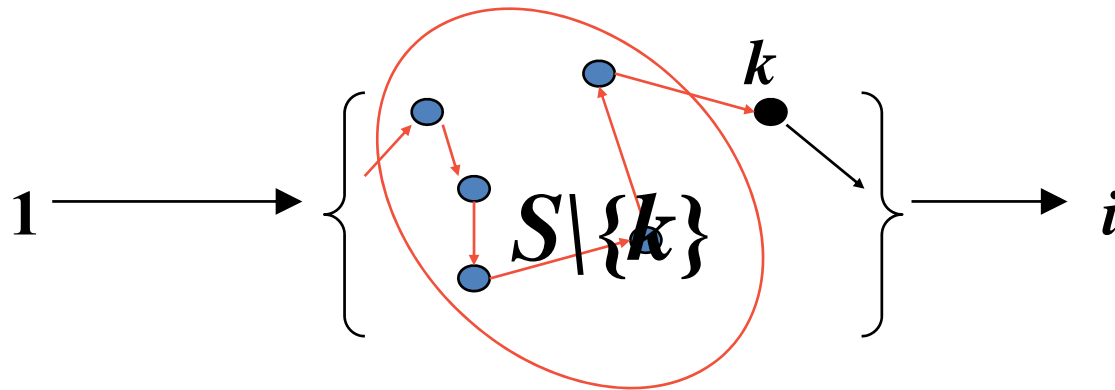
Then the TSP is equivalent to the problem of finding $\mathbf{cost}[\{2,3,\dots,n\},1]$



TSP: Dynamic Programming

The lengths $\mathbf{cost}[S,i]$ have to satisfy the following equations:

$$\mathbf{cost}[S,i] = \min_{k \in S} \{ \mathbf{cost}[S \setminus \{k\}, k] + c_{ki} \}$$

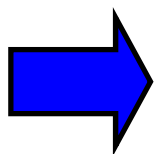


THE TRAVELLING SALESMAN PROBLEM: DP

How many entries does the **cost**[S,i] table have?

$$\begin{aligned}
 & (n-1) * (n-2) + \binom{n-1}{2} * (n-3) + \binom{n-1}{3} * (n-4) + \binom{n-1}{4} * (n-5) + \dots + \binom{n-1}{n-2} * 1 + 1 \\
 & = 2 * \binom{n-1}{2} + 3 * \binom{n-1}{3} + 4 * \binom{n-1}{4} + 5 * \binom{n-1}{5} + \dots + (n-1) * \binom{n-1}{n-1} + 1 \\
 & = (n-1) * 2^{n-2} - (n-2)
 \end{aligned}$$

n=	5	10	20	30	40	50	100	1000
TSP	29	2296	4980718	7784628196	1.072E+13	1.3792E+16	3.1374E+31	2.67E+303
Knaps B=500	2500	5000	10000	15000	20000	25000	50000	500000



Computational not a feasible approach for larger number of cities